

Recursivias & Aplicaciones

- Torre de Hanoi
- Disk Pile Problem
- Programacion Dinamica

Torre de Hanoi

```
Mover n, A, B, C) {  
  if (n == 1) Mover 1, noDis(A, C);  
  else if (n > 1) {
```

```
    Mover (n-1, A, C, B)
```

```
    Mover (1, A, B, C)
```

```
    Mover (n-1, B, A, C)
```

```
  }
```

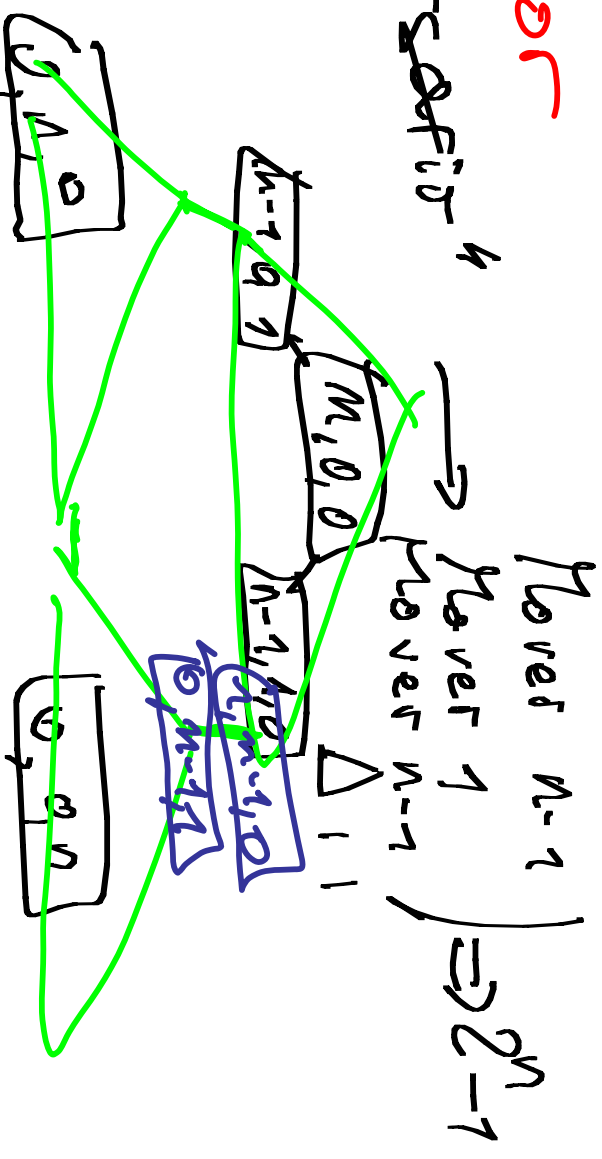
Completed

- CoTa Superior

$$f(n) = \# \text{ movimientos de Mover}(n, A, B, C) \\ = 2 f(n-1) + 1 \\ = 2^n - 1 \quad \quad \quad 1 \quad \quad 2 \left(2^{n-1} - 1 \right) + 1$$

- CoTa Inferior

$\hookrightarrow$ "adversario"
 $\hookrightarrow$ grafica



II Disk Pile Problem

n_i discos de tamanho i

$i \in [1..n]$, $n = \sum n_i$ discos

Complexidade

por n fixo

- peor caso $2^n - 1$

- caso superior: $2^n - 1$

- caso inferior: $2^n - 1$

por n, h fixos

- peor caso

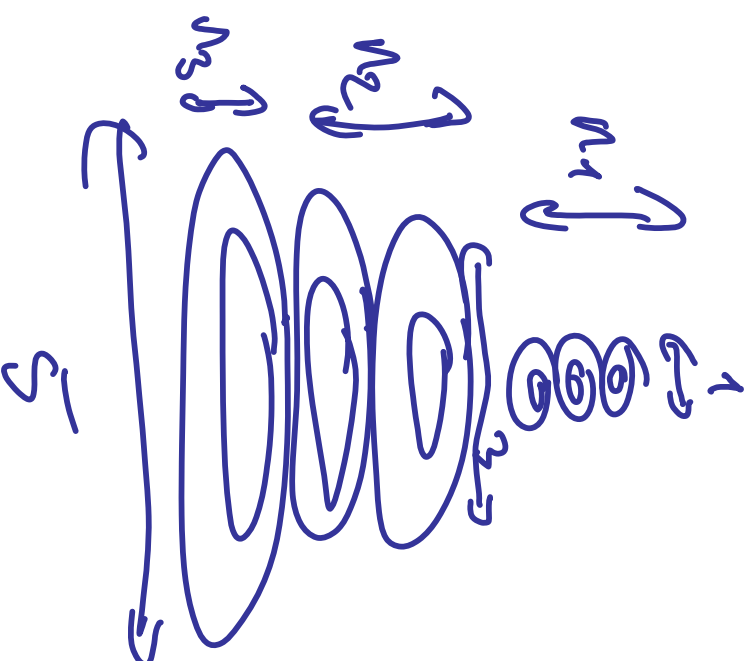
$$n_1 = n - h + 1 \Rightarrow 2^{n-h+1} - 1$$

$$n_1 = 1 \Rightarrow 2^{n-1} - 1$$

- peor caso por $n_1, n_2, \dots, n_h$ fixos

$$f(n_1, n_2, \dots, n_h) = 2^{f(n_1-1, n_2, \dots, n_h)} + n_h =$$

$$\sum_{i=1}^h 2^{n_i}$$



III Recurrences

$$\left(\sum_{i=1}^n 2^{h-i} \right) / n/h = \frac{n}{h} \left(\sum_{i=0}^{h-1} 2^i \right)$$

$$= \frac{n}{h} (2^h - 1)$$

$$\boxed{\sum_{i=1}^n i} = \frac{n(n+1)}{2}$$

$$\boxed{\sum_{i=1}^n q^i} = \frac{1-q^{n+1}}{1-q}$$

$$= \text{primero termino} \frac{1 - \text{razon}^{\text{num term}}}{1 - \text{razon}}$$

$$\boxed{\sum_{i=1}^n i^2} = ?$$

Recurrences Lineals

Fibonacci

$$F_0 = 1$$

$$F_1 = 1$$

$$F_n = F_{n-2} + F_{n-1}$$

$$\rightarrow \lambda^n = \lambda^{n-2} + \lambda^{n-1}$$

$$\rightarrow \lambda^n - \lambda^{n-2} - \lambda^{n-1} = 0$$

$$\lambda^{n-2}(\lambda^2 - \lambda - 1) = 0$$

$$\begin{cases} \lambda = 0 \\ \lambda^2 - \lambda - 1 = 0 \end{cases}$$

$$\Rightarrow \lambda = \frac{1 \pm \sqrt{5}}{2}$$

General $F_n = a_1 F_{n-1} + a_2 F_{n-2} + \dots + a_k F_{n-k} + a_0$

$$\begin{pmatrix} F_{n-1} \\ F_{n-2} \\ \vdots \\ F_{n-k+1} \end{pmatrix} = \begin{pmatrix} a_1 & a_2 & \dots & a_k & a_0 \\ 1 & 0 & \dots & 0 & 0 \\ 0 & 1 & \dots & 0 & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \dots & 1 & 0 \end{pmatrix}$$

$$\begin{pmatrix} F_{n-1} \\ \vdots \\ F_{n-k} \end{pmatrix} = \begin{pmatrix} F_k \\ \vdots \\ F_1 \end{pmatrix} A_{n-k+1}$$

