

**AS4501-CC4501**

# **Complejidad Computacional**

## **Básica**

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# Notacion

Introduction to algorithms, Thomas H. Cormen ... [et al.]. 3rd ed.

$\Theta(g(n)) = \{f(n) : \text{there exist positive constants } c_1, c_2, \text{ and } n_0 \text{ such that } 0 \leq c_1g(n) \leq f(n) \leq c_2g(n) \text{ for all } n \geq n_0\}$ .<sup>1</sup>

$O(g(n)) = \{f(n) : \text{there exist positive constants } c \text{ and } n_0 \text{ such that } 0 \leq f(n) \leq cg(n) \text{ for all } n \geq n_0\}$ .

$\Omega(g(n)) = \{f(n) : \text{there exist positive constants } c \text{ and } n_0 \text{ such that } 0 \leq cg(n) \leq f(n) \text{ for all } n \geq n_0\}$ .

<sup>1</sup>Within set notation, a colon means “such that.”

# Ejemplos

- Multiplicar un vector de  $n$  elementos por un escalar  $\mathcal{O}(n)$
- Multiplicar 2 vectores de  $n$  elementos  $\mathcal{O}(n)$
- Multiplicar una matriz de  $n \times m$  por un vector de  $m$  elementos  $\mathcal{O}(nm)$
- Encontrar un elemento en un arreglo ordenado  $\mathcal{O}(\log n)$
- DFT 
$$X_k = \sum_{n=0}^{N-1} x_n e^{\frac{-2\pi i}{N} kn}$$
  $\mathcal{O}(n^2)$
- FFT  $\mathcal{O}(n \log n)$
- Convolución 1D vecotres de tamaño  $n$  y  $m$   $\mathcal{O}(nm)$
- Convolucion 2D  $\mathcal{O}(n_x n_y m^2)$
- Convolucion 2D separable en  $x$  e  $y$   $\mathcal{O}(n_x n_y m)$   
$$\mathcal{O}(n_x n_y (\log n_x + \log n_y) + (n_x + n_y)m \log m)$$