

# AMPL and CPLEX tutorial

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## 1 The steel production problem

### 1.1 The problem

2 products can be produced at a steel mill:

- We can make 200 tons of product 1 in an hour; the profit for each ton is 25 dollars; the demand is 6000 tons. We must make at least 1000 tons of this product.
- We can make 140 tons of product 2 in an hour; the profit for each ton is 30 dollars; the demand is 4000 tons. We must make at least 2000 tons of this product.

We have 40 hours of production time available.

The goal is to design a production plan to maximize total profit.

With  $x_i =$  tons of product  $i$  to be made, we get the following LP:

$$\begin{aligned} \max \quad & 25x_1 && +30x_2 \\ \text{st.} \quad & x_1 \geq 0, && x_2 \geq 0 \\ & \frac{1}{200}x_1 && + \frac{1}{140}x_2 \leq 40 \\ & 1000 \leq x_1 \leq 6000 \\ & 2000 \leq x_2 \leq 4000 \end{aligned} \tag{1.1}$$

### 1.2 Writing and running a correct model

The simplest version of the steel problem's solution is below:

```
### File: steel-simple.mod
```

```

var x1 >=1000, <= 6000;
var x2 >=2000, <= 4000;

maximize profit: 25*x1 + 30*x2;

subject to time: (1/200)*x1 + (1/140)*x2 <= 40;

## here: 'profit' and 'time'
## are arbitrarily chosen names.

```

-----

AMPL run (the output may look slightly different, depending on what version you are using; in particular, in the student version, that you get from [ampl.com](http://ampl.com), you need to type ‘‘option solver cplex;’’, and in the version on the department’s Unix machines, you need to type ‘‘option solver cplexamp;’)

```

ampl: model steel.mod;
ampl: option solver cplex;
ampl: solve;
ILOG CPLEX 8.000, licensed to "university-chapel hill, nc", options: e m b q
CPLEX 8.0.0: optimal solution; objective 188571.4286
0 dual simplex iterations (0 in phase I)
ampl: display x1, x2;
x1 = 1000
x2 = 4000
ampl:

```

When we split the problem into model and data files, they look like this:

```

### File: steel.mod

param n :=2;

param a {j in 1..n};
param b;
param c {j in 1..n};
param u {j in 1..n}; # Upper bound on production
param l {j in 1..n}; # Lower bound on production

var x {j in 1..n} <= u[j], >= l[j];

maximize profit: sum {j in 1..n} c[j] * x[j];

```

```
subject to time: sum {j in 1..n} (1/a[j]) * x[j] <= b;
```

```
## here: 'param' and 'var' are reserved keywords; 'profit' and 'time'  
## are arbitrarily chosen names.
```

```
-----  
### File: steel.dat
```

```
param      a  1 200  
           2 140;
```

```
param      c  1 25  
           2 30;
```

```
param      u  1 6000  
           2 4000;
```

```
param      l  1 1000  
           2 2000;
```

```
param b := 40;
```

```
-----  
AMPL run (the output may look slightly different, depending on  
what version you are using).
```

```
ampl: model steel.mod;  
ampl: data steel.dat;  
ampl: option solver cplexamp;  
ampl: solve;  
ILOG CPLEX 8.000, licensed to "university-chapel hill, nc", options: e m b q  
CPLEX 8.0.0: optimal solution; objective 188571.4286  
0 dual simplex iterations (0 in phase I)  
ampl: display x;  
x [*] :=  
1 5142.86  
2 2000  
;  
  
ampl:
```

```
### Remark: the number of iterations can be 0,
```

```
### if the problem is very simple;
### in that case, a so-called ‘‘LP preprocessor’’
### already solves the problem.
```

```
### We can change some of the data, and resolve:
```

```
ampl: let u[1] := 5000;
ampl: display x.ub;
x.ub [*] :=
1 5000
2 4000
;
```

```
### x.ub is the current upper bound on x; the .ub extension works for
### any variable. I.e. we can write y.ub, if y is a variable vector.
### Similarly, x.lb gives the lower bounds.
```

```
ampl: solve;
ILOG CPLEX 8.000, licensed to "university-chapel hill, nc", options: e m b q
CPLEX 8.0.0: optimal solution; objective 188000
1 dual simplex iterations (0 in phase I)
```

### 1.3 Inputting a general problem

This is how to input an LP of the form

$$\begin{aligned} \max \quad & c^T x \\ \text{st.} \quad & Ax \leq b \end{aligned}$$

where  $A$  is a matrix with  $m$  rows, and  $n$  columns.

The  $A, b, c$  for the problem will be

$$A = \begin{pmatrix} 5 & 6 \\ 7 & 8 \\ 9 & 10 \end{pmatrix}, b = \begin{pmatrix} 13 \\ 15 \\ 19 \end{pmatrix}, c = \begin{pmatrix} 5 \\ 6 \end{pmatrix}.$$

Model file:

```
param n;
```

```

param m;

param A {i in 1..m, j in 1..n};
param b {i in 1..m};
param c {j in 1..n};

var x {j in 1..n};

maximize whatever: sum {j in 1..n} c[j] * x[j];

subject to cons {i in 1..m}:
    sum {j in 1..n} A[i,j]*x[j] <= b[i];

```

Data file:

```

param n :=2;
param m :=3;

param A:   1 2 :=
           1 5 6
           2 7 8
           3 9 10;

param b :=
           1 13
           2 15
           3 19;

param c :=
           1 5
           2 6;

```

## 1.4 Debugging an incorrect model

Suppose we have set up the data in a way, so the LP is infeasible. Usually in an infeasible LP, there are only a few constraints which result in infeasibility. As an extreme example, in:

$$x_1 \leq 2, x_1 \geq 3, 0 \leq x_i \leq 1 (i = 2, \dots, 1000)$$

the only constraints that cause trouble, are the first two. They form a so called *irreducible infeasible system*; that is, a subset of all inequalities, which are infeasible, but

dropping any one of them would make this system feasible. (i.e.  $x_1 \leq 2$ ,  $x_1 \geq 3$  is infeasible, but dropping any one of these gives just one inequality, which is of course feasible).

For instance, this data makes the steel problem infeasible:

```
### File: steel.dat
```

```
param      a  1 200  
           2 140;
```

```
param      c  1 25  
           2 30;
```

```
param      u  1 6000  
           2 4000;
```

```
param      l  1 4000  
           2 3000;
```

```
param b := 40;
```

```
AMPL run:
```

```
ampl: solve;
```

```
presolve: constraint time cannot hold:
```

```
body <= 40 cannot be >= 41.4286; difference = -1.42857
```

```
### Not too useful info... We will find an IIS, to localize the problem.
```

```
ampl: option presolve 0;
```

```
### This tells the solver to turn the preprocessor off.
```

```
ampl: option cplex_options 'iisfind 1';
```

```
### This tells the solver to find an IIS.
```

```
ampl: solve;
```

```
ILOG CPLEX 8.000, licensed to "university-chapel hill, nc", options: e m b q
```

```

CPLEX 8.0.0: iisfind 1
Bound infeasibility column 'x1'.
CPLEX 8.0.0: infeasible problem.
0 simplex iterations (0 in phase I)
Returning iis of 2 variables and 1 constraints.
constraint.dunbdd returned
1 extra dual simplex iterations for ray (1 in phase I)

```

```
suffix iis symbolic OUT;
```

```

option iis_table '\
0      non      not in the iis\
1      low      at lower bound\
2      fix      fixed\
3      upp      at upper bound\
';
suffix dunbdd OUT;

```

```

### Most of the above stuff is just technicalities
### that you can ignore...
### The important part comes below:

```

```

ampl: display x.iis;
x.iis [*] :=
1 low
2 low
;

```

```

ampl: display time.iis;
time.iis = upp

```

The meaning of the above lines is:

$$x_1 \geq 4000, x_2 \geq 3000, (1/200)x_1 + (1/140)x_2 \leq 40$$

is an IIS. (That is, the upper bounds on  $x$  have nothing to do with the infeasibility).

## 1.5 Another variant of the steel problem

The next model is the same, but it gives names to the products.

```
### File: steel2.mod
```

```
set P;
```

```
param a {j in P};
```

```
param b;
```

```
param c {j in P};
```

```
param u {j in P};
```

```
param l {j in P};
```

```
var x {j in P};
```

```
maximize profit: sum {j in P} c[j] * x[j];
```

```
subject to time: sum {j in P} (1/a[j]) * x[j] <= b;
```

```
subject to limit {j in P}: l[j] <= x[j] <= u[j];
```

```
### File: steel2.dat
```

```
set P := bands coils;
```

```
param:  a :=  bands 200  
       coils 140;
```

```
param:  c :=  bands 25  
       coils 30;
```

```
param:  u :=  bands 6000  
       coils 4000;
```

```
param:  l :=  bands 1000  
       coils 2000;
```

```
param b := 40;
```

```
### AMPL run:
```

```
ampl: model steel2.mod;
```

```
ampl: data steel2.dat;
```

```
ampl: option solver cplexamp;
```

```
ampl: solve;
```

```
ILOG CPLEX 8.000, licensed to "university-chapel hill, nc", options: e m b q
CPLEX 8.0.0: optimal solution; objective 188571.4286
0 dual simplex iterations (0 in phase I)
ampl: display x;
x [*] :=
bands 5142.86
coils 2000
;
```

**Important!!** If you make a mistake in a model, or data file, you will need to 1) fix it, 2) type “reset”, or “reset data” before you can reread those files. Example:

```
ampl: model steel.mod;

steel.mod, line 11 (offset 205):
    syntax error
context: m >>> aximize <<< profit: sum {j in 1..n} c[j] * x[j];
ampl?

### There was a mistake in the model file; we fix it, and reread it.

ampl: reset;
ampl: model steel.mod;
ampl: data steel.dat;

steel.dat, line 1 (offset 2):
    syntax error
context: p >>> aram <<< a 1 200
ampl?

### Now the model file was OK, but there is a mistake in the data file;
### we fix the data file, and reread only the data file (the model file
### was OK to start with).

ampl? ;
ampl: reset data;
ampl: data steel.dat;
ampl:
```

## 2 The minimum cost flow problem

This problem is excellent to illustrate how to define variables  $x_{ij}$ , where the *existing* variables are just a small subset of the *possible* ones.

```
### File: mcf.mod

param n :=5;      # Number of nodes;

set ARCS within {1..n, 1..n};

param demand {1..n};
  check: sum {i in 1..n} demand[i] = 0;

### This statement will check that the sum of demands is zero, as one would
### expect for the problem to be feasible.

param cost {ARCS};
param u {ARCS};

var x {ARCS} >=0;

minimize total_cost:
  sum { (i,j) in ARCS } cost[i,j]*x[i,j];

subject to balance {i in 1..n}:
  sum { (j,i) in ARCS } x[j,i] - sum{ (i,j) in ARCS } x[i,j] = demand[i];

### File: mcf.dat

param demand := 1 1
2 3
3 5
4 -6
5 -3;

param: ARCS: cost := 1 2 10
  1 4 5
                1 5 7

  2 3 5
  2 4 6
  2 1 1
  3 1 5
```

```

3 4 10
3 5 1
4 2 1
4 5 6
5 1 1
5 2 3
5 4 7;

```

### 3 A multiperiod problem

Suppose we have a multiperiod problem, with variables

- $inv_i$  (for inventory at the end of period  $i$ ),  $prod_i$  (for production in period  $i$ , and
- parameters  $demand_i$  (for demand in period  $i$ ),
- constraints

$$inv_{i-1} + prod_i = demand_i + inv_i$$

We can write these constraints concisely as follows (only parts of the model and data files are written down):

```

### File: production.mod

param n;

param demand {1..n};

var inv {0..n}; # inventory;
var prod {1..n}; # production;

subject to balance {i in 1..n}:
    inv[i-1] + prod[i] = demand[i] + inv[i];

etc.

### File: production.dat

param n := 6;

param demand := 1 5
2 3
3 16

```

```
4 11
5 10
        6 7;
```

etc.

This has the following advantages, as opposed to writing out the 6 constraints individually:

- This is much cleaner, and easier to read.
- If the parameter  $n$  is not “hardwired” into the program, i.e. you do not write 6 in any place where  $n$  should be used, then the code is much more flexible. If you have 2 data files, one with say  $n = 6$ , the other with  $n = 1000$ , then you can use the same model file with both.

## 4 Some neat tricks

There are some useful internal variables in AMPL:

```
_nvar is the number of variables;
_var   is a vector containing the values of all variables;
_varname is a vector containing the names of all variables;
```

So in the steel problem, we can do:

```
AMPL: display _varname;
_varname [*] :=
1  'x[1]'
2  'x[2]'
;
```

```
AMPL: display _var;
_var [*] :=
1  5142.86
2  2000
;
```

```
AMPL: display {i in 1.._nvars: _var[i]>0} _varname[i], {i in 1.._nvars: _var[i]>0} _var[i];
```

```
: _varname[i]  _var[i]    :=  
1  'x[1]'      5142.86  
2  'x[2]'      2000  
;
```

The last command displays the names and values of all variables with a positive value (in this instance, actually all variables have a positive value).